

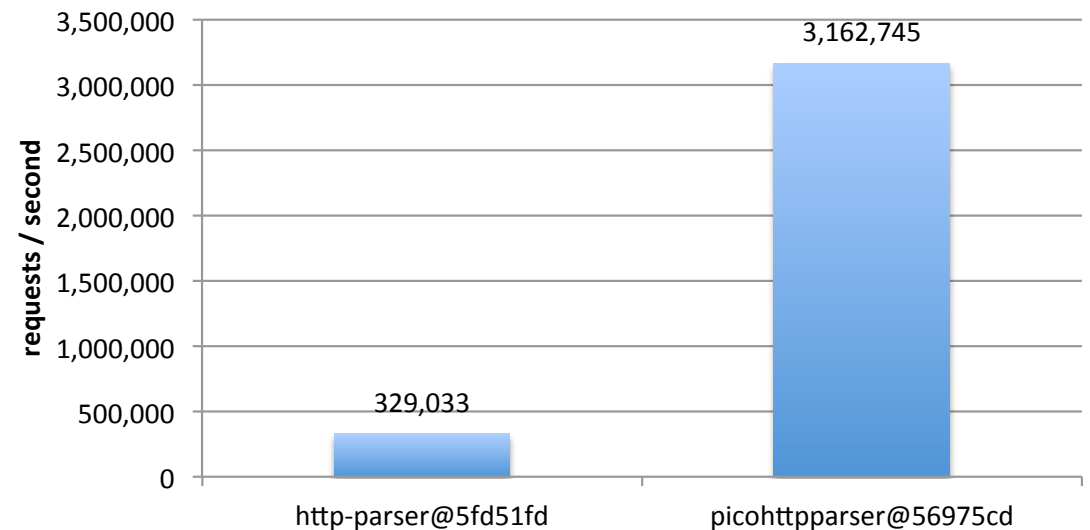
# Writing a Fast HTTP Parser

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# Parsing input

- HTTP/1 request parser may or may not be a bottleneck, depending on its performance
  - if the parser is capable of handling 1M reqs/sec, then it will spend 10% of time if the server handles 100K reqs/sec.

HTTP/1 Parser Performance Comparison (2014)



# How fast could a text parser be?

- around 1GB/sec. is a good target
  - since any parser needs to read every byte and execute a conditional branch depending on the value
    - # of instructions: 1 load + 1 inc + 1 test + 1 conditional branch
    - would likely take several CPU cycles (even if superscalar)
    - unless we use SIMD instructions

# Parsing input

## ■ What's wrong with this parser?

```
for (; s != end; ++s) {  
    int ch = *s;  
    switch (ctx.state) {  
    case AAA:  
        if (ch == ' ' )  
            ctx.state = BBB;  
        break;  
    case BBB:  
        ...  
    }  
}
```

# Parsing input (cont'd)

- never write a character-level state machine if performance matters

```
for (; s != end; ++s) {  
    int ch = *s;  
    switch (ctx.state) { // ← executed for every char  
    case AAA:  
        if (ch == ' ' )  
            ctx.state = BBB;  
        break;  
    case BBB:  
        ...  
    }  
}
```

# Parsing input fast

- each state should consume a sequence of bytes

```
while (s != end) {
    switch (ctx.state) {
    case AAA:
        do {
            if (*s++ == ' ') {
                ctx.state = BBB;
                break;
            }
        } while (s != end);
        break;
    case BBB:
        ...
    }
```

# Stateless parsing

- stateless in the sense that no *state* value exists
  - stateless parsers are generally faster than stateful parsers, since it does not have *state* - a variable used for a conditional branch
- HTTP/1 parsing can be stateless since the request-line and the headers arrive in a single packet (in most cases)
  - and even if they did not, it is easy to check if the end-of-headers has arrived (by looking for CR-LF-CR-LF) and then parse the input
    - this countermeasure is essential to handle the Slowloris attack

# pichttpparser is stateless

- states are the execution contexts (instead of being a variable)

```
const char* parse_request(const char* buf, const char* buf_end, ...)
{
    /* parse request line */
    ADVANCE_TOKEN(*method, *method_len);
    ++buf;
    ADVANCE_TOKEN(*path, *path_len);
    ++buf;
    if ((buf = parse_http_version(buf, buf_end, minor_version, ret)) == NULL)
        return NULL;
    EXPECT_CHAR('\015');
    EXPECT_CHAR('\012');
    return parse_headers(buf, buf_end, headers, num_headers, max_headers, ...);
}
```



# loop exists within a function (≠state)

- the code looks for the end of the header value

```
#define IS_PRINTABLE(c) ((unsigned char)(c) - 040u < 0137u)

static const char* get_token_to_eol(const char* buf, const char* buf_end, ...
{
    while (likely(buf_end - buf >= 8)) {
#define DOIT() if (unlikely(! IS_PRINTABLE(*buf))) goto NonPrintable; ++buf
        DOIT(); DOIT(); DOIT(); DOIT();
        DOIT(); DOIT(); DOIT(); DOIT();
#undef DOIT
        continue;
    NonPrintable:
        if ((likely((uchar)*buf < '\040') && likely(*buf != '\011'))
            || unlikely(*buf == '\177'))
            goto FOUND_CTL;
    }
}
```

# The hottest loop of picohttpparser (cont'd)

- after compilation, uses 4 instructions per char

```
movzbl  (%r9), %r11d
movl    %r11d, %eax
addl    $-32, %eax
cmpl    $94, %eax
ja      LBB5_5
movzbl  1(%r9), %r11d    // load char
leal    -32(%r11), %eax  // subtract
cmpl    $94, %eax       // and check if is printable
ja      LBB5_4           // if not, break
movzbl  2(%r9), %r11d    // load next char
leal    -32(%r11), %eax  // subtract
cmpl    $94, %eax       // and check if is printable
ja      LBB5_15          // if not, break
movzbl  3(%r9), %r11d    // load next char
```

...

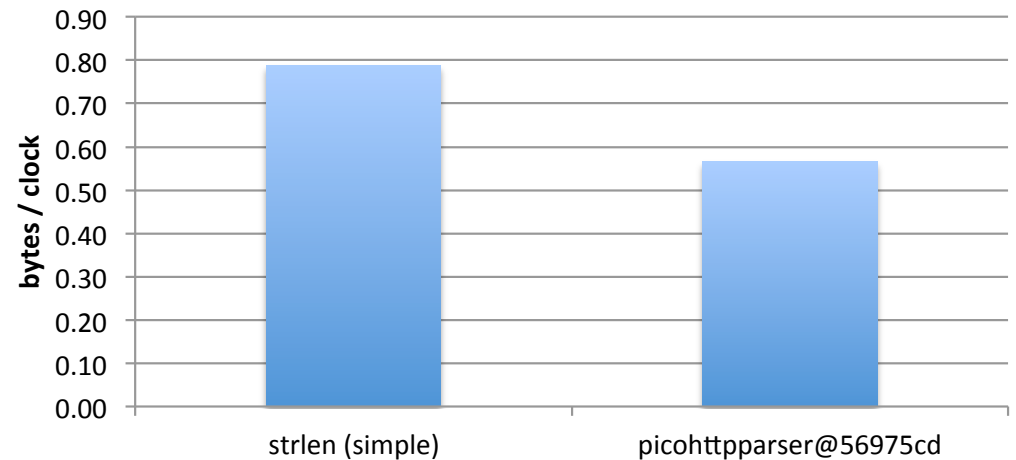
# strlen vs. picohttpparser

- not as fast as strlen, but close

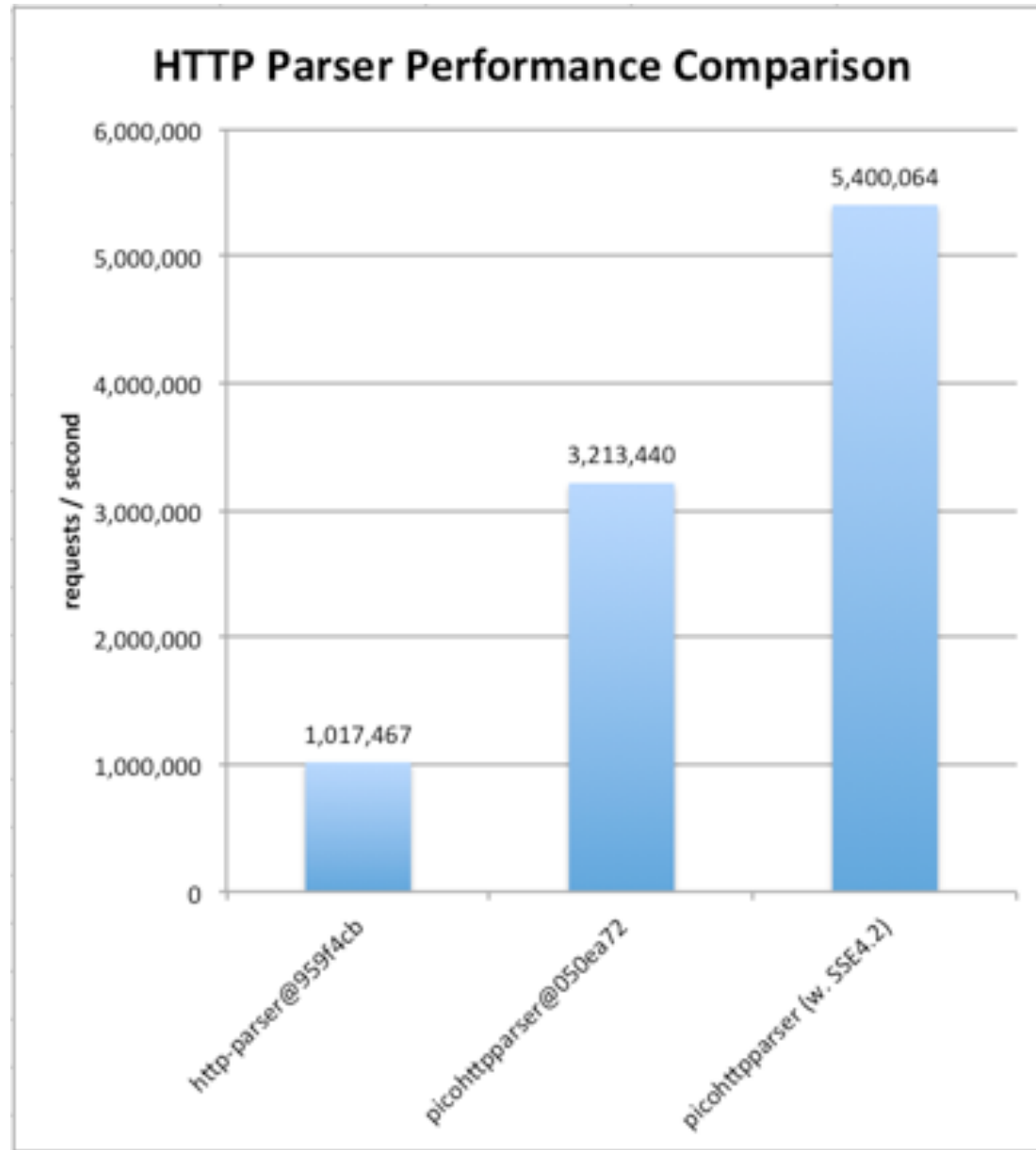
```
size_t strlen(const char *s) {  
    const char *p = s;  
    for (; *p != '\0'; ++p)  
        ;  
    return p - s;  
}
```

- not much room left for further optimization (wo. using SIMD insns.)

strlen vs. picohttpparser (2014)



# in 2015, we've did SIMD, got 2x speed



latest version uses SSE 4.2 insns; [a fork that runs 2x faster using AVX2 insns](#). also exists (developed by CloudFlare)

# picohttpparser is small and simple

```
$ wc picohttpparser.?
```

```
  376    1376   10900 picohttpparser.c
```

```
   62     333    2225 picohttpparser.h
```

```
  438    1709   13125 total
```

```
$
```

- good example of do-it-simple-for-speed approach

# conclusion

- text-based parser can be pretty fast
  - so that the overhead compared to binary-based approach is negligible
- we should not introduce binary-based structure for encoding headers in the future versions of HTTP, in the notion that it would be faster